

## INTRODUCTION

Unconstrained optimization problems have been paid considerable attention by the researchers, because of comprehensive practical application background. There are many authors have made great efforts on the study of optimization algorithms. Consider a general unconstrained optimization problem denoted by

$$
\begin{equation*}
\min _{x \in R^{n}} f(x) \tag{1}
\end{equation*}
$$

Where $f(x)$ is a continuously differentiable function from $R^{n}$ to $R$. In the literature, it is customary to use iterative methods to solve this problem. At current iteration $x_{k}$, if $g_{k}=\nabla f\left(x_{k}\right) \neq 0$, we can find a step-length $\alpha_{k}$ by carrying some line search along the direction $d_{k}$, and then obtain the next iteration as

$$
\begin{equation*}
x_{k+1}=x_{k}+\alpha_{k} d_{k} \tag{2}
\end{equation*}
$$

where $d_{k}$ is the search direction, which can be determined by many methods. The literature ${ }^{[1]}$ points out that the Broyden-Fletcher-Goldfarb-Shanno (BFGS) algorithm is the best quasi-Newton method and gave the associated procedures. For convenience, we will employ the BFGS algorithm in this paper to determine $d_{k}$, namely, let
$d_{k}=\left\{\begin{array}{cc}-H_{k} g_{k}, & \text { if } g_{k}^{T} d_{k}<0 \\ -g_{k}, & \text { if } g_{k}^{T} d_{k} \geq 0\end{array}\right.$

Where $g_{k}=\nabla f\left(x_{k}\right) \neq 0, H_{k}$ is generated by the BFGS correction formula. Moreover, (3) guarantee that the condition $-g_{k}^{T} d_{k}>0$ holds.

There are already some well-known rules used to determine the step-length $\alpha_{k}$ in (2). Among them, the Arimijo rule, the Goldstein rule, and the Wolfe rule are popular and widely used by many authors ${ }^{[1]}$. However, in obtaining the step length $\alpha_{k}$, traditional line searches require the function value to decrease monotonically at every iteration, namely

$$
\begin{equation*}
f\left(x_{k+1}\right) \leq f\left(x_{k}\right) \tag{4}
\end{equation*}
$$

Consequently, they are in general called to be monotone line search technique, which is resultful in some situations. However, subsequent studies showed that the convergence rate of monotone line search technique may reduce considerably when the iteration locates in a narrow curved valley ${ }^{[2,3,4]}$. To overcome this problem, Grippo et al. introduced a highly innovative method ${ }^{[2]}$ called the nonmonotone line search techniqueand illustrated its effectiveness by means of some numerical tests. This method has been developed by many authors. For those nonmonotone line search rules ${ }^{[2-13]}$, the inequality $f\left(x_{k+1}\right)>f\left(x_{k}\right)$ may hold for any $k$, and therefore, it can play a nonmonotone search role in the above three rules. However, the nonmonotone line search rules ${ }^{[2-9,11-13]}$ for problems (1) essentially required the approximation sequence $\left\{x_{k}\right\}$ satisfies $f\left(x_{k}\right) \leq f\left(x_{0}\right)$ for any $k \geq 1$. Under this condition, the approximation sequence $\left\{x_{k}\right\}$ will be trapped and can not escape from the valley bottom if the initial point $x_{0}$ locates near a valley bottom. Motivated by this problem, we propose a new rule called to be a nonmonotone line search combination rule in the following. Moreover, we show that it possesses the global convergence property by virtue of ${ }^{[3]}$ and ${ }^{[4]}$. With the help of numerical experiments it is shown that the proposed method is very effective for above problem.

## Nonmonotone line search combination rule

In this section, we will put forward a nonmonotone line search combination rule. First, the following assumption is the necessary.

Assumption 1 Throughout this paper, we assume that the function $f: R^{n} \rightarrow R$ is bounded and differentiable on the level set $\Omega=\left\{x\left|x \in R^{n}: f(x) \leq c\right| f\left(x_{0}\right) \mid\right\}$ for a given constant $c \geq 1$.

Nonmonotone line search combination rule. Let the bounded step-length $\alpha_{k} \geq 0$ along the direction $d_{k}$ such that

$$
\begin{equation*}
f\left(x_{k+1}\right)=f\left(x_{k}+\alpha_{k} d_{k}\right) \leq \sum_{r=o}^{m(k)} \lambda_{k r} \beta^{h_{k r} \operatorname{sign} f\left(x_{k-r}\right)} f\left(x_{k-r}\right)+\rho \alpha_{k} g_{k}^{T} d_{k} \tag{5}
\end{equation*}
$$

Where $\beta \geq 1, h_{k r} \geq 0$, and $m(k)=\min [k, M-1]$ with $M \geq 1$ is a positive integer. Let $1 \geq \lambda_{k r}>\lambda>0, \sum_{r=0}^{m(k)} \lambda_{k r}=1$, $h_{k}=\sum_{r=0}^{m(k)} h_{k r}$ and $\sum_{k=0}^{\infty} h_{k}=S$ with $S$ is a finite constant.

Clearly, inserting $M=1$ and $\beta=1$ into the representation (5) gives $f\left(x_{k}+\alpha_{k} d_{k}\right) \leq f\left(x_{k}\right)+\rho \alpha_{k} g_{k}^{T} d_{k}$, then the rule (5) reduces to the rule of Arimijo. Further, if $f\left(x_{k-l}\right)=\max _{0 \leq r \leq m(k)}\left\{f\left(x_{k-r}\right)\right\}$ and $\lambda_{k l}=1, \lambda_{k r}=0(r \neq l)$, then (5) becomes $f\left(x_{k}+\alpha_{k} d_{k}\right) \leq \max _{0 \leq j \leq m(k)}\left\{f\left(x_{k-j}\right)\right\}+\rho \alpha_{k} g_{k}^{T} d_{k}$, which is the nonmonotone line search rule proposed by Grippo et al. ${ }^{[2]}$. Therefore, the new rule (5) is a general version of some traditional rules well known in the literature.

## Proof for the global convergence of the new rule

In this section, we investigate the global strong convergence properties of unconstrained optimization in conjunction with the new nonmonotone line search combination rule. First, the following some definitions and lemmas are necessary.

Definition 1 The function $\sigma:[0,+\infty) \rightarrow\left[0,+\infty\right.$ ) is a forcing function (F-function) if for any sequence $\left\{t_{i}\right\} \subset[0,+\infty$ ), $\lim _{i \rightarrow \infty} \sigma\left(t_{i}\right)=0$ implies $\lim _{i \rightarrow \infty} t_{i}=0$.

The literature ${ }^{[3]}$ proved that there exists a F-function $\sigma\left(t_{i}\right)$ such that
$\rho \alpha_{k} g_{k}^{T} d_{k} \leq-\sigma\left(t_{k}\right)$ with $t_{k}=-g_{k}^{T} d_{k} /\left\|d_{k}\right\| \geq 0$.
By the inequalities (5) and (6), the rule (5) becomes

$$
\begin{equation*}
f\left(x_{k+1}\right)=f\left(x_{k}+\alpha_{k} d_{k}\right) \leq \sum_{r=0}^{m(k)} \lambda_{k r} \beta^{h_{k r} \operatorname{sig} n f\left(x_{k-r}\right)} f\left(x_{k-r}\right)-\sigma\left(t_{k}\right) . \tag{7}
\end{equation*}
$$

Lemma 3.1 In (2), the search direction $d_{k}$ is determined by the following BFGS algorithm. Then there is a constant $\gamma>0$ satisfying

$$
\begin{equation*}
-g_{k}^{T} d_{k} /\left\|d_{k}\right\| \geq \gamma\left\|g_{k}\right\|, k=0,1,2, \cdots \tag{8}
\end{equation*}
$$

Proof. When the $g_{k} \neq 0$, according (3), we easily get $g_{k}^{T} d_{k}<0$. In view of $-g_{k}^{T} d_{k} /\left\|d_{k}\right\|=\left\|g_{k}\right\| \cos \left(-g_{k}, d_{k}\right)$, then there is a constant $\gamma>0$ such that $\cos \left(-g_{k}, d_{k}\right)>\gamma>0$. Therefore, (8) holds.

Lemma 3.2 If $\alpha_{k}$ satisfies the rule (5) for $k \geq 1$, then
$f\left(x_{k+1}\right) \leq\left|f\left(x_{0}\right)\right| \beta^{\sum_{=0}^{k} n_{i}}-\lambda \sum_{r=0}^{k-1} \sigma\left(t_{r}\right)-\sigma\left(t_{k}\right)$
Proof. The principle of mathematical induction will be used to prove the conclusion.
If $k=1$ and $M=1$, then we have $m(k)=0$. By the inequality (7), we have
$f\left(x_{2}\right) \leq \lambda_{10} \beta^{h_{10} \operatorname{sign} f\left(x_{1}\right)} f\left(x_{1}\right)-\sigma\left(t_{1}\right) \leq \beta^{h_{10}} f\left(x_{1}\right)-\sigma\left(t_{1}\right)$.
Noting that $\beta^{h_{k}} \geq \beta^{h_{k r}} \geq 1 \geq h_{k r}>\lambda$ and
$f\left(x_{1}\right) \leq \lambda_{00} \beta^{h_{00} \operatorname{sign} f\left(x_{0}\right)} f\left(x_{0}\right)-\sigma\left(t_{0}\right)=\beta^{h_{00}}\left|f\left(x_{0}\right)\right|-\sigma\left(t_{0}\right)$.
We have

$$
\begin{aligned}
& f\left(x_{2}\right) \leq \beta^{h_{10}}\left\{\beta^{h_{00}}\left|f\left(x_{0}\right)\right|-\sigma\left(t_{0}\right)\right\}-\sigma\left(t_{1}\right) \\
& =\beta^{h_{00}+h_{01}}\left|f\left(x_{0}\right)\right|-\beta^{h_{10}} \sigma\left(t_{0}\right)-\sigma\left(t_{1}\right)
\end{aligned}
$$

$$
\begin{equation*}
\leq \beta^{h_{0}+h_{1}}\left|f\left(x_{0}\right)\right|-\beta^{h_{10}} \sigma\left(t_{0}\right)-\sigma\left(t_{1}\right) \tag{10}
\end{equation*}
$$

If $k=1$ and $M \geq 2$, by the inequality (7), we have
$f\left(x_{2}\right) \leq \lambda_{10} \beta^{h_{10} \operatorname{sign} f\left(x_{1}\right)} f\left(x_{1}\right)+\lambda_{11} \beta^{h_{11} \operatorname{sign} f\left(x_{0}\right)} f\left(x_{0}\right)-\sigma\left(t_{1}\right)$
$\leq \lambda_{10} \beta^{h_{10}}\left\{\beta^{h_{00}}\left|f\left(x_{0}\right)\right|-\sigma\left(t_{0}\right)\right\}+\lambda_{11} \beta^{h_{11}}\left|f\left(x_{0}\right)\right|-\sigma\left(t_{1}\right)$
$\leq\left(\lambda_{10} \beta^{h_{1}+h_{0}}+\lambda_{11} \beta^{h_{1}}\right)\left|f\left(x_{0}\right)\right|-\lambda_{10} \beta^{h_{1}} \sigma\left(t_{0}\right)-\sigma\left(t_{1}\right)$
$\leq \beta^{h_{1}+h_{0}}\left(\lambda_{10}+\lambda_{11}\right)\left|f\left(x_{0}\right)\right|-\lambda_{10} \sigma\left(t_{0}\right)-\sigma\left(t_{1}\right)$
$\leq \beta^{h_{1}+h_{0}}\left|f\left(x_{0}\right)\right|-\lambda \sigma\left(t_{0}\right)-\sigma\left(t_{1}\right)$.

By (10) and (11), it follows that (9) holds for $k=1$. We now assume that (9) holds for $j=k-1$, namely,
$f\left(x_{k}\right) \leq\left|f\left(x_{0}\right)\right| \beta^{\sum_{i=0}^{k-1} h_{i}}-\lambda \sum_{r=o}^{k-2} \sigma\left(t_{r}\right)-\sigma\left(t_{k-1}\right)$.
Then we have

$$
\begin{aligned}
& f\left(x_{k+1}\right) \leq \sum_{r=o}^{m(k)} \lambda_{k r} \beta^{h_{k r} s i g n f\left(x_{k-r}\right)} f\left(x_{k-r}\right)-\sigma\left(t_{k}\right) \\
& \leq \sum_{r=o}^{m(k)} \lambda_{k r} \beta^{h_{k}}\left\{\left|f\left(x_{0}\right)\right| \beta^{\sum_{i=0}^{k-r-1} h_{i}}-\lambda \sum_{i=o}^{k-r-2} \sigma\left(t_{i}\right)-\sigma\left(t_{k-r-1}\right)\right\}-\sigma\left(t_{k}\right) \\
& \leq \sum_{r=o}^{m(k)} \lambda_{k r} \beta^{h_{k}}\left\{\left|f\left(x_{0}\right)\right| \beta^{\sum_{i=0}^{k-1} h_{i}}-\lambda \sum_{i=o}^{k-r-2} \sigma\left(t_{i}\right)-\sigma\left(t_{k-r-1}\right)\right\}-\sigma\left(t_{k}\right) \\
& \leq \sum_{r=o}^{m(k)} \lambda_{k r} \beta^{h_{k}}\left\{\left|f\left(x_{0}\right)\right| \beta^{\sum_{i=0}^{k-1} h_{i}}-\lambda \sum_{i=o}^{k-m(k)-2} \sigma\left(t_{i}\right)-\sigma\left(t_{k-r-1}\right)\right\}-\sigma\left(t_{k}\right) \\
& \leq\left|f\left(x_{0}\right)\right|\left(\sum_{r=o}^{m(k)} \lambda_{k r}\right) \beta^{\sum_{i=0}^{k} h_{i}}-\lambda\left(\sum_{r=o}^{m(k)} \lambda_{k r}\right) \sum_{i=o}^{k-m(k)-2} \sigma\left(t_{i}\right)-\sum_{r=o}^{m(k)} \lambda_{k r} \sigma\left(t_{k-r-1}\right)-\sigma\left(t_{k}\right) \\
& =\left|f\left(x_{0}\right)\right| \beta^{\sum_{i=0}^{k} h_{i}}-\lambda \sum_{i=o}^{k-m(k)-2} \sigma\left(t_{i}\right)-\sum_{r=o}^{m(k)} \lambda_{k r} \sigma\left(t_{k-r-1}\right)-\sigma\left(t_{k}\right) \\
& \leq\left|f\left(x_{0}\right)\right| \beta^{\sum_{i=0}^{k} h_{i}}-\lambda \sum_{i=o}^{k-m(k)-2} \sigma\left(t_{i}\right)-\lambda \sum_{r=o}^{m(k)} \sigma\left(t_{k-r-1}\right)-\sigma\left(t_{k}\right) \\
& =\left|f\left(x_{0}\right)\right| \beta^{\sum_{i=0}^{k} h_{i}}-\lambda \sum_{i=o}^{k-m(k)-2} \sigma\left(t_{i}\right)-\lambda \sum_{i=k-m(k)-1}^{k-1} \sigma\left(t_{i}\right)-\sigma\left(t_{k}\right) \\
& \leq\left|f\left(x_{0}\right)\right| \beta^{\sum_{i=0}^{k} h_{i}}-\lambda \sum_{r=o}^{k-1} \sigma\left(t_{r}\right)-\sigma\left(t_{k}\right) .
\end{aligned}
$$

Which means that (9) holds for $j=k$. By the principle of mathematical induction, (9) holds for any given $k \geq 1$.
Next, we will prove the global strong convergence of the new rule.
Theorem 3.1 Under the above assumptions 1, let the search direction $d_{k}$ and the step-length $\alpha_{k}$ be determined by BFGS algorithm and (5), respectively. Assume $\left\{x_{k}\right\}$ is a sequence generated by (2) according to the search direction $d_{k}$ and the step-length $\alpha_{k}$, Then we have
$\left\{x_{k}\right\} \in \Omega$, and $\lim _{k \rightarrow \infty}\left\|g_{k}\right\|=0$.
$f\left(x_{k+1}\right) \leq\left|f\left(x_{0}\right)\right| \beta^{\sum_{i=1}^{k}}-\lambda \sum_{r=o}^{k-1} \sigma\left(t_{r}\right)-\sigma\left(t_{k}\right) \leq\left|f\left(x_{0}\right)\right| \beta^{S}$.
Therefore, $f\left(x_{k+1}\right) \leq c\left|f\left(x_{0}\right)\right|$ with $c=\beta^{s}>1$. Then, by the definition of $\Omega,\left\{x_{k}\right\} \in \Omega$ follows immediately. By (9), we have

$$
\begin{aligned}
& f\left(x_{k+1}\right) \leq\left|f\left(x_{0}\right)\right| \beta^{\sum_{i=0}^{k} h_{i}}-\lambda \sum_{r=o}^{k-1} \sigma\left(t_{r}\right)-\sigma\left(t_{k}\right) \\
& \leq\left|f\left(x_{0}\right)\right| \beta^{\sum_{i=0}^{k} h_{i}}-\lambda \sum_{r=o}^{k} \sigma\left(t_{r}\right) \\
& \leq c\left|f\left(x_{0}\right)\right|-\lambda \sum_{r=o}^{k} \sigma\left(t_{r}\right) .
\end{aligned}
$$

Which implies
$0 \leq \lambda \sum_{r=o}^{k} \sigma\left(t_{r}\right) \leq c\left|f\left(x_{0}\right)\right|-f\left(x_{k+1}\right)$.
for all $k \geq 1$. By Assumption $1, f\left(x_{k}\right)$ is bounded on the level set $\Omega$. Then (12) indicates that $\lambda \sum_{r=0}^{k} \sigma\left(t_{r}\right)<\infty$ which implies $\lim _{k \rightarrow \infty} \sigma\left(t_{k}\right)=0$. Combined with the fact that function $\sigma$ is a F-function, it follows then that $\lim _{k \rightarrow \infty} t_{k}=\lim _{k \rightarrow \infty}\left(-g_{k}^{T} d_{k} /\left\|d_{k}\right\|\right)=0$.

By Lemma3.1, it follows that $\lim _{k \rightarrow+\infty}\left\|g_{k}\right\|=0$. The proof is completed.

## Numerical tests

Here, we apply our rule to some standard tests problems.
Algorithm (I)
Step 1. Initialization. Given the initial values $x_{0} \in R^{n}, H_{0} \in I$ and other data including an integer $M \geq 1$, a constant $\alpha_{0}=1, \varepsilon \geq 0, \rho \in(0,0.5)$, as well as $k=0$.

Step 2. Test termination conditions. Examine the stopping criterion by computing $g_{k}=\nabla f\left(x_{k}\right)$. If $\left\|g_{k}\right\| \leq \varepsilon, x^{*}=x_{k}$ and the algorithm stops.

Step 3. Determine search direction. Calculate $d_{k}=-H_{k} g_{k}$, if $g_{k}^{T} d_{k}>0$, set $d_{k}=-g_{k}$, which guarantee the condition $-g_{k}^{T} d_{k}>0$ holds.

Step 4. Determine the line search step $\alpha_{k}$. Let $m(k)=\min [k, M-1]$. If (5) holds, $\alpha_{k}=\alpha$. Otherwise, contract $\alpha$.
Step 5. Compute the next point. Set $s_{k}=\alpha_{k} d_{k}$ and $x_{k+1}=x_{k}+s_{k}$, and then calculate $f\left(x_{k+1}\right)$ and $g_{k+1}=\nabla f\left(x_{k+1}\right)$.
Step 6. Update the iteration matrix $H_{k+1}$ using BFGS formulae, namely, Set
$H_{k+1}=H_{k}+\frac{s_{k} s_{k}^{T}}{s_{k}^{T} y_{k}}-\frac{H_{k} y_{k} y_{k}^{T} H_{k}}{y_{k}^{T} H_{k} y_{k}}+v_{k} v_{k}^{T}$
With $y_{k}=g_{k+1}-g_{k}, v_{k}=\left(y_{k}^{T} H_{k} y_{k}\right)^{\frac{1}{2}}\left[\frac{s_{k}}{s_{k}^{T} y_{k}}-\frac{H_{k} y_{k}}{y_{k}^{T} H_{k} y_{k}}\right]$. Set $k=\mathrm{k}+1$ and then go to the step 2.
The functions in the numerical tests are the same as ${ }^{[5,7]}$. Take now the parameters involved in Algorithm (I) as follows:

$$
\varepsilon=10^{-6}, \rho=10^{-3}, H_{0}=I_{n \times n}, \lambda_{k r}=\frac{1}{1+m(k)}, h_{k r}=\frac{1}{(1+\mathrm{k})^{1.2}} .
$$

With $r=0,1, \cdots, m(k)$.
Clearly, Algorithm (I) reduces to the standard of the BFGS algorithm (saying Algorithm (II)) when $M=1$ and $\beta=1$. If $M>1$ and $\beta=1$, algorithm (I) is similar to the nonmonotone BFGS algorithm (saying Algorithm (III)). The three methods are compared with the aid of Matlab, where the controllable parameter $M$ is taken as 1 and 6 while $\lambda$ as 1 and 3 . The numerical results are presented in the following four TABLES, in which $n_{g}$ and $n_{f}$ denote the outside loop iterations and the function evaluations, respectively, while $f\left(x^{*}\right)$ is the function value of the approximate solution $x^{*}$.

TABLE 1 : Algorithm (I) does not dominate compared to Algorithm (II) or (III)

| Problem | Dimension | Algorithm 2( $M=1, \beta=1$ ) |  | Algorithm 3( $M=3, \beta=1$ ) |  | Algorithm 1( $M=\mathbf{3 , ~} \beta=\mathbf{6}$ ) |  |
| :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: |
|  |  | $n_{g} / n_{f}$ | $f\left(x^{*}\right)$ | $n_{g} / n_{f}$ | $f\left(x^{*}\right)$ | $n_{g} / n_{f}$ | $f\left(x^{*}\right)$ |
|  |  |  |  |  | ----- |  |  |
|  |  |  | ----- |  | $9.0191 \mathrm{e}-$ |  |  |
|  |  |  | 2.8656e- |  | , |  |  |
|  |  |  | 13 |  | 7.5025e- |  |  |
| Pow.Sin. |  |  | $4.0157 \mathrm{e}-$ |  | 13 |  |  |
| Ex.Ros. | 4 | >999 | 13 | >999 | 4.6651e- | 60/327 | $3.3350 \mathrm{e}-04$ |
| Ex.Ros. | 10 | 95/245 | 3.1755e- | 98/246 | 20 | 85/212 | $4.6259 \mathrm{e}-15$ |
| Ex.Ros. | 30 | 190/515 | 13 | >999 | 3.5514e- | 175/464 | 6.1892e-14 |
| Ex.W.\& | 40 | 230/636 | 4.6651e- | 239/649 | 23 | 205/560 | $4.0964 \mathrm{e}-14$ |
| H. | 2 | 32/77 | 20 | 32/77 | $2.5945 \mathrm{e}-$ | 28/66 | $1.2043 \mathrm{e}-17$ |
| Ge.Ros. | 2 | 39/96 | $1.5647 \mathrm{e}-$ | 38/92 | 16 | 31/75 | 7.8019e-20 |
| Ge.Ros. | 8 | >999 | 17 | 60/167 | $1.2228 \mathrm{e}-$ | 66/180 | 1.8386e-17 |
| Ge.Ros. | 20 | >999 | ----- | 133/379 | 17 | 138/385 | $7.5640 \mathrm{e}-17$ |
| Penalty I | 30 | 136/328 | ----- | 22/60 | $2.4773 \mathrm{e}-$ | 22/60 | $2.4773 \mathrm{e}-04$ |
| Penalty I | 40 | >999 | 2.4773e- | >999 | 04 | 74/160 | $3.3925 \mathrm{e}-04$ |
| Penalty I | 80 | $>999$ | 04 | $>999$ | ----- | 155/351 | 7.1305e-04 |
| Penalty II | 32 | >999 | ----- | 104/248 | ------ | 112/270 | 0.1032 |
| Penalty II | 40 | >999 | ------ | 98/253 | 0.1032 | 102/253 | 0.5569 |
| Penalty II Watson | 10 | 66/154 | ------ | 68/155 | 0.5569 | 64/149 | $5.3212 \mathrm{e}-07$ |
| Watson Watson | 36 | 58/157 | ------ | 85/210 | 5.3203e- | 56/150 | $3.3854 \mathrm{e}-08$ |
| Watson <br> Watson | 150 | >9999 | $5.3203 \mathrm{e}-$ | 97/356 | ${ }_{0}^{07}$ | 98/344 | $1.2336 \mathrm{e}-09$ |
| Watson | 70 | >999 | $07$ | 90/349 | $1.5352 \mathrm{e}-$ | 90/347 | 5.8369e-15 |
| Per.Qua | 80 | >999 | $3.3840 \mathrm{e}-$ | 103/403 | 09 | 102/399 | $1.6813 \mathrm{e}-14$ |
| Per.Qua Per.Qua | 90 | >999 | 08 | >999 | 1.2814e- | 115/451 | 1.0610e-14 |
| Per.Qua | 80 | >999 | --- | 201/3828 | 09 | 103/217 | 324.00 |
| Raydan I | 100 | 519/14606 | ----- | 73/192 | 1.0450e- | 146/318 | 505.00 |
| Raydan I | 200 | >999 | ----- | >999 | 14 | 297/752 | $2.0100 \mathrm{e}+03$ |
| Raydan I |  |  | ------- |  | $\begin{gathered} 2.9824 \mathrm{e}- \\ 15 \end{gathered}$ |  |  |
|  |  |  | 505.00 |  | ----- |  |  |
|  |  |  | ----- |  | 324.00 |  |  |
|  |  |  |  |  | 505.00 |  |  |
|  |  |  |  |  | ----- |  |  |

Form TABLE 1, Algorithm (I) performs the best for the given initial point. As see form TABLE 2, the efficiency of Algorithm (I) is not worse then Algorithm (II) and Algorithm (III).

Extended Freudenstein \& Roth (EFR) function is non-convex with two minimum points, one of which is globally minimum. Form TABLE 3, it is concluded that Algorithm (I) performs the best in finding the globally optimal solution and is very suitable for solving high-dimension problem. In specification, in the case of the dimension equal to 2, both Algorithms (II) and (III) give the minimum 48.9843 while Algorithm (I) gives the globally minimum 2.0835e-019.

The results for Brown and Dennis Function are presented in TABLE 4. It is easy to see that the iterations of Algorithm (I) is decreasing as $\$$ becomes larger, which indicates the efficiency the new algorithm.

TABLE 2 : Algorithm (I) does not dominate compared to Algorithm (II) or (III)

| Problem | Dimension | Algorithm 2( $M=1, \beta=1)$ |  |  | Algorithm 3( $M=\mathbf{3}, \beta=1)$ |  | $\begin{gathered} \text { Algorithm } 1(M=3, \beta \\ =6) \end{gathered}$ |
| :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: |
|  |  | $n_{g} / n_{f}$ | $f\left(x^{*}\right)$ | $n_{g} / n_{f}$ | $f\left(x^{*}\right)$ | $n_{g} / n_{f}$ | $f\left(x^{*}\right)$ |
| Beale |  |  |  |  |  |  |  |
| Ex.W.\& |  |  |  | 2/28 |  | 2/28 |  |
| H. | 2 | 12/28 | $9.6837 \mathrm{e}-18$ | 89/225 | $9.6837 \mathrm{e}-18$ | 83/206 | $9.6837 \mathrm{e}-18$ |
| Ex.W.\& | 10 | 85/219 | $6.5620 \mathrm{e}-16$ | 126/34 | 3.7749e-16 | 121/32 | $1.8803 \mathrm{e}-14$ |
| H. | 20 | 123/341 | $1.4192 \mathrm{e}-14$ | 1 | $4.7872 \mathrm{e}-14$ | 4 | $6.2400 \mathrm{e}-15$ |
| Ex.W.\& | 30 | 156/448 | 3.6872e-14 | 154/44 | 6.5794e-14 | 153/43 | $1.7713 \mathrm{e}-14$ |
| H. | 1000 | 22/45 | $1.2011 \mathrm{e}-06$ | 0 | $1.2011 \mathrm{e}-06$ | 2 | $1.2011 \mathrm{e}-06$ |
| BDEXP | 5000 | 23/47 | 2.8855e-06 | 22/45 | 2.8855e-06 | 22/45 | 2.8855e-06 |
| BDEXP | 2 | 8/20 | $8.0664 \mathrm{e}-07$ | 23/47 | 8.0664e-07 | 23/47 | 8.0664e-07 |
| Penalty II | 10 | 18/51 | $4.5920 \mathrm{e}-14$ | 9/21 | $3.9012 \mathrm{e}-14$ | 9/21 | $6.1450 \mathrm{e}-14$ |
| Per.Qua | 100 | 11/65 | $100$ | 19/51 | $100$ | 19/50 | $100$ |
| Raydan II | 1000 50 | 11/65 | 1000 54517 | 16/78 | 1000 | 17/78 | 1000 |
| Raydan II | 500 | 53/109 | 4.4742e-07 | 44/89 | 4.4742e-07 | 44/89 | 4.4742e-07 |
| Trigo. | 1000 | 12/36 | 883.1941 | 52/107 | 883.1941 | 52/107 | 883.1941 |
| Trigo. | 100 | 75/219 | $1.1267 \mathrm{e}-08$ | 16/45 | 1.1268e-08 | 21/55 | $1.1307 \mathrm{e}-08$ |
| Ex.Penalty <br> Watson |  |  |  | $76 / 217$ |  | $76 / 216$ |  |

TABLE 3 : Results for Extended Freudenstein and Roth function (using the initial point $x_{0}=(0.5,-2,0.5,-2, \cdots, 0.5,-2)$ )

| Dimension | Algorithm 2 $(M=\mathbf{1 , ~} \beta=\mathbf{1})$ |  | Algorithm 3(M=3, $\beta=\mathbf{1})$ |  | Algorithm 1 $(M=\mathbf{3}, \beta=\mathbf{6})$ |  |
| :---: | :---: | :---: | :---: | :---: | :---: | :---: |
|  | $n_{g} / n_{f}$ | $f\left(x^{*}\right)$ | $n_{g} / n_{f}$ | $f\left(x^{*}\right)$ | $n_{g} / n_{f}$ | $f\left(x^{*}\right)$ |
| 2 | $10 / 32$ | 48.9843 | $11 / 33$ | 48.9843 | $15 / 42$ | $2.0835 \mathrm{e}-19$ |
| 6 | $45 / 759$ | 146.9528 | $22 / 65$ | 146.9528 | $39 / 158$ | $1.1415 \mathrm{e}-15$ |
| 10 | $27 / 95$ | 244.9213 | $30 / 95$ | 244.9213 | $46 / 144$ | $1.3625 \mathrm{e}-16$ |
| 18 | $53 / 174$ | 440.8583 | $50 / 167$ | 440.8583 | $62 / 217$ | $2.8598 \mathrm{e}-16$ |
| 22 | $52 / 219$ | 538.8268 | $60 / 198$ | 538.8268 | $75 / 259$ | $1.7857 \mathrm{e}-16$ |
| 24 | $46 / 504$ | 587.8110 | $62 / 206$ | 587.8110 | $80 / 282$ | $1.6609 \mathrm{e}-16$ |

TABLE 4 : Results for Brown and Dennis function (using $\mathbf{m}=10$ and the initial point $x_{0}=(0.5,-\mathbf{2}, \mathbf{0 . 5},-\mathbf{2}, \cdots, 0.5,-\mathbf{2})$ )

| $\begin{gathered} \text { Algorithm2 } \\ (M=1, \beta=1) \end{gathered}$ | $\begin{aligned} & \text { Algorithm } 3 \\ & (M=3, \beta=1) \end{aligned}$ | Algorithm 1( $M=3$ ) |  |  |  |  |  |
| :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: |
|  |  | $\beta=6$ | $\beta=10$ | $\beta=20$ | $\beta=30$ | $\beta=40$ | $\beta=50$ |
| $n_{g} / n_{f}$ | $n_{g} / n_{f}$ | $n_{g} / n_{f}$ | $n_{g} / n_{f}$ | $n_{g} / n_{f}$ | $n_{g} / n_{f}$ | $n_{g} / n_{f}$ | $n_{g} / n_{f}$ |
| ----- | ----- | 134/609 | 78/380 | 37/154 | 37/150 | 35/145 | 31/127 |

## CONCLUSIONS

By (12), we get $c\left|f\left(x_{0}\right)\right|-f\left(x_{k}\right) \geq 0$. Then it follows that $f\left(x_{k}\right) \leq c\left|f\left(x_{0}\right)\right|$, where $c \geq 1$, and therefore it is possible to ensure $f\left(x_{k}\right)>f\left(x_{0}\right)$. Hence, we can reach the goal of slackness and let the iteration point escape from the valley near $x_{0}$ and search a better solution. This idea is a breakthrough of this paper.

We mention here that the rule (5) can be easily achieved. In general, there are many selections for $\lambda_{k r}, h_{k r}$ which is involved in (5). Take $h_{k r}=\frac{1}{(1+\mathrm{k})^{p}}$ with the constants $p>1$ for example. In this case, the series p-series $\sum_{k=0}^{\infty} \frac{1}{(k+1)^{p}}$ converge on a limited number $S$.

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